Aims:

- *Resource Sharing*: costly resources (high-quality printers and expensive programs) can be shared
- *Computation Speedup*: Algorithms can run concurrently
- *Reliability*: Failure of one site does not imply failure of the whole system
  Redundancy prerequisite

Especially last point difficult to fulfill

Several Levels of Distribution possible
listed from tightly coupled to loosely coupled

- Shared memory
- Shared file system
- Bus Systems
- Switching Systems

Key Problem Areas:

- Transparency: pretend to be one computer
- Reliability: want availability, fault tolerance
- Performance
- Scalability: avoid centralised tables and algorithms

Simplest way of communication:
Client/Server Model

Idea: group of processes (servers) used by clients
Advantage: Simple communication

Simple enough to study several problems
First problem: Addressing! Possible Solutions:

- Hardware address into client code: inflexible
- Broadcasting: works only on local networks
- Name Servers: Ask special host
  Example: Domain Name Service, DNS

Second Problem:
Blocking (synchronous) vs. non-blocking primitives
Conceptual ease vs. performance

Third Problem:
Reliable vs. unreliable primitives
Where does the error correction go?
Kernel (once for all) vs. application (possibly more efficient)

Remote Procedure Call

Very simple idea: execute procedure on different host
Goal: total transparency

Basic Schema:

- Client sends arguments to server
- Server executes call
- Server sends results back

Difference to local call hidden in kernel routines

Details complicated:
- Have to transfer parameters
- deal with failures

Problems with parameter passing

- Different representation on different machines: either common format (inefficient), or store format in message
- What to do with call-by-reference parameters? Can copy arrays, but not arbitrary data structures

Failure problems more complicated
Have several cases

- Client cannot locate server: Generate exception
  ⇒ transparency lost
- Lost Request Messages: use timer
- Lost Reply Messages: client timer insufficient:
  could execute operation more than once
  Solution: use sequence numbers
- Server Crashes: Sequence numbers not enough: When did crash occur?
  Can guarantee *at least once*, `at most once` semantics, but not *exactly once* semantics ⇒ have to have call-specific remedies
- Client Crash: leaves orphans (unwanted computations)
  No general way of getting rid of them

Generating Timestamps

Unique timestamps needed for co-ordination
No problem in monoprocessor system: use system clock
Not possible in distributed systems
One way out:

- Each host maintains logical clock which is advanced with each event
- All message from host contain logical clock
- When message with greater logical clock is received, increase own logical clock to this value
- with two identical timestamps, let host number decide

A Distributed Version (Ricart and Agrawala)
Assume reliable messages, unique timestamps
Following steps:

- Process trying to enter critical section: sends to all other processes name of section and unique timestamp

- Process receiving such a message:
  - Sends back OK if not interested in critical section
  - Queues message if already in critical section
  - Receiver wants to enter critical section
    ⇒ enters critical section if its request has lower timestamp and queues message, otherwise sends OK

Grants mutual exclusion without deadlock or starvation

Problems:

- Requires that everyone knows about all other hosts
- Algorithm fails if one host fails
  ⇒ Reconstruction of network necessary in this case

Suitable for networks where configuration is stable

A Token Ring Algorithm

Assumption: Network organised on a (physical or logical) ring, *i.e.*,
each node has unique successor in line
Simple algorithm:

- At initialisation, generate token
- Pass token around continuously
- Process wanting to enter a critical region waits for token
- After leaving critical section, process passes token to next neighbour

Properties:

- Detecting lost tokens difficult: Time spent in critical region unbound
- Detecting dead processes easy if sent token is acknowledged

Election Algorithms

Problem: Select new co-ordinator
Assumption:
Know id of *every* host on the network
First example: Bully algorithm

- P sends message to all hosts with higher number
- No response ⇒ P wins and is co-ordinator
- Answer received ⇒ host with higher number has taken over

Second Example: Ring Algorithm

- any host may send *Election* message
- passed around the net, which each hostid added
- If original host gets message, determines co-ordinator and sends new message around
- After it has gone round, host removes it