



Dependency Quantified Boolean Formulas

Uwe Bubeck

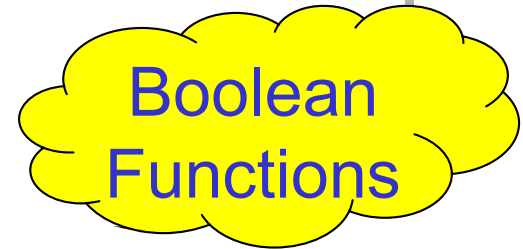
International Graduate School
Dynamic Intelligent Systems,
University of Paderborn

27.3.2007

QBF and Beyond



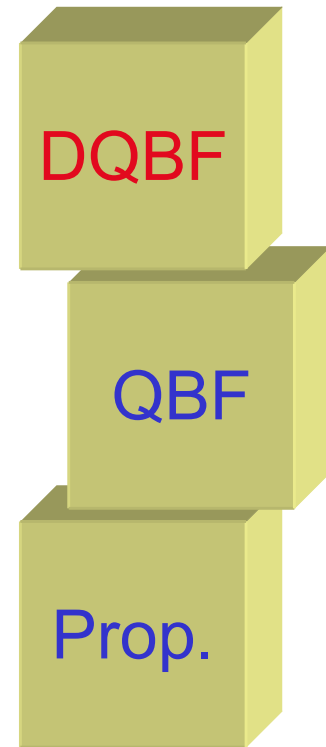
Quantified Boolean Formulas (QBF)
extend propositional logic with
quantifiers over atoms.



Further extension:
quantifiers with explicit dependencies
→ **Dependency QBF (DQBF)**

Motivations:

- short representations
- elegant/natural modeling

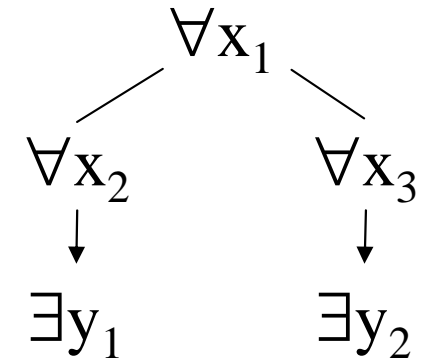


Partially-Ordered Quantifiers 1/2



Consider a quantified formula where:

- y_1 depends on x_1 and x_2
- y_2 depends on x_1 and x_3



Problem:

By definition, QBFs are in **prenex form** with a **linear order of quantifiers**.

→ we must **flatten** the **quantifier hierarchy**

$$\forall x_1 \forall x_2 \exists y_1 \forall x_3 \exists y_2 \phi(x_1, x_2, y_1) \wedge \psi(x_1, x_3, y_2)$$

now: y_2 depends on x_1, x_2, x_3

→ **loss of information, larger search space**



Solution:

Partially-ordered Quantification

[Henkin '61; Peterson et al. '79]

→ Dependency Quantified Boolean Formulas

$$\forall x_1 \forall x_2 \forall x_3 \exists y_1(x_1, x_2) \exists y_2(x_1, x_3) \phi(x_1, x_2, y_1) \wedge \psi(x_1, x_3, y_2)$$

Ability to indicate for each existential variable on which universals it depends.

Advanced Dependencies



No strict correspondence between dependencies and formula structure required.

Example:

$$\forall x_1 \forall x_2 \forall x_3 \exists y_1(x_1, x_2) \exists y_2(x_1, x_3)$$

$$\phi(x_1, x_2, y_1) \wedge \psi(x_1, x_3, y_2) \wedge \tau(y_1, y_2)$$

not possible in
(non-prenex) QBF!

Why combine „unrelated“ variables in one subformula?
allows to impose additional restrictions on the choice of existential variables.

→ increases expressive power

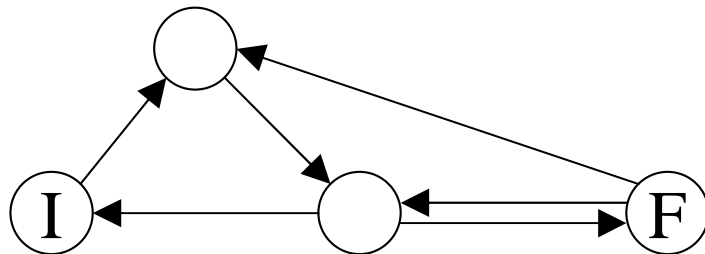
→ new modeling approaches



DQBF Modeling

How can we leverage the **increased expressiveness** of DQBF for **reducing formula sizes**?

Modeling Example: **Graph Reachability**



Is there a path of length k from an initial state to a final state?

→ DQBF modeling as **three-player game**.

→ Only **$O(\log k)$ variables** vs. $O(k)$ in QBF.



DQBF Complexity

DQBF-SAT is known to be NEXPTIME-complete (vs. PSPACE-completeness of QBF).

Can we identify easier/tractable subclasses?

Dependency Quantified Horn Formulas:

→ can be solved in time $O(|\forall| \cdot |\Phi|)$ (same as QHORN).

→ can be transformed into equivalent QHORN formulas of length $O(|\forall| \cdot |\Phi|)$.